

# SAT-based Verification of NSPKT Protocol Including Delays in the Network

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# Importance of Security Protocols

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- Key point of security systems
- Used in many areas
- Errors in: structure, operations, security
- Specification and verification importance
- Need for the complete formal model
- IT market development sets new requirements

# New Challenges

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- Detailed analysis of the protocol
- "Tailor-made" security
- The importance of time

# World Leaders

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- AVISPA
- ProVerif
- Scyther
- Verlcs
- PRISM

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# Basic elements

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Protocol definition (with time aspect)

$$\begin{aligned}\alpha^1 &= (S_{\rightarrow}, R_{\leftarrow}, L), \\ \alpha^2 &= (\tau, D, X, G, tc)\end{aligned}\tag{1}$$

# Example: Needham-Schroeder Public Key Protocol

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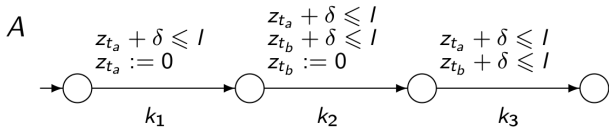
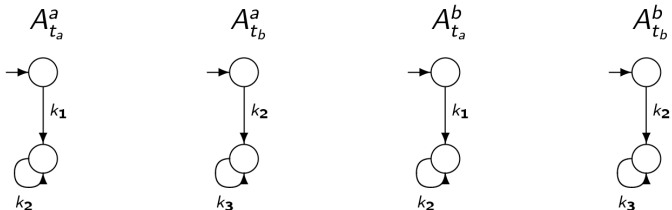
$$\begin{aligned}\alpha_1 &= (\alpha_1^1, \alpha_1^2), \\ \alpha_1^1 &= (\mathcal{A}; \mathcal{B}; \langle \tau_{\mathcal{A}} \cdot \mathcal{I}_{\mathcal{A}} \rangle_{\mathcal{K}_{\mathcal{B}}}), \\ \alpha_1^2 &= (\tau_1; \mathcal{D}_1; \{\tau_{\mathcal{A}}, \mathcal{I}_{\mathcal{A}}, \mathcal{K}_{\mathcal{B}}\}; \{\tau_{\mathcal{A}}\}; \tau_1 + \mathcal{D}_1 - \tau_{\mathcal{A}} \leq \mathcal{L}),\end{aligned}$$

$$\begin{aligned}\alpha_2 &= (\alpha_2^1, \alpha_2^2), \\ \alpha_2^1 &= (\mathcal{B}; \mathcal{A}; \langle \tau_{\mathcal{B}} \cdot \tau_{\mathcal{A}} \rangle_{\mathcal{K}_{\mathcal{A}}}), \\ \alpha_2^2 &= (\tau_2; \mathcal{D}_2; \{\tau_{\mathcal{B}}, \tau_{\mathcal{A}}, \mathcal{K}_{\mathcal{A}}\}; \{\tau_{\mathcal{B}}\}; \tau_2 + \mathcal{D}_2 - \tau_{\mathcal{A}} \leq \mathcal{L} \wedge \tau_2 + \mathcal{D}_2 - \tau_{\mathcal{B}} \leq \mathcal{L}),\end{aligned}$$

$$\begin{aligned}\alpha_3 &= (\alpha_3^1, \alpha_3^2), \\ \alpha_3^1 &= (\mathcal{A}; \mathcal{B}; \langle \tau_{\mathcal{B}} \rangle_{\mathcal{K}_{\mathcal{B}}}), \\ \alpha_3^2 &= (\tau_3; \mathcal{D}_3; \{\tau_{\mathcal{A}}, \mathcal{K}_{\mathcal{B}}\}; \{\emptyset\}; \tau_3 + \mathcal{D}_3 - \tau_{\mathcal{A}} \leq \mathcal{L} \wedge \tau_3 + \mathcal{D}_3 - \tau_{\mathcal{B}} \leq \mathcal{L}).\end{aligned}$$

# Automata Model

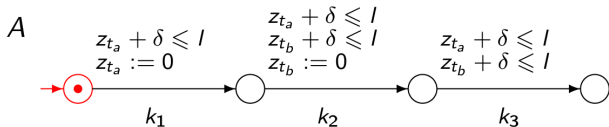
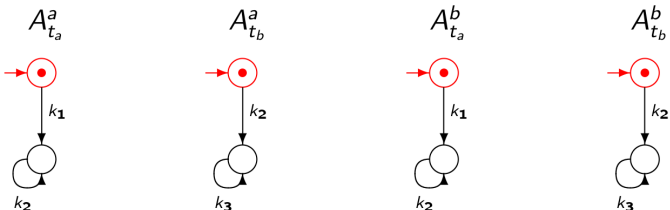
Network of synchronized timed automata for NSPK protocol





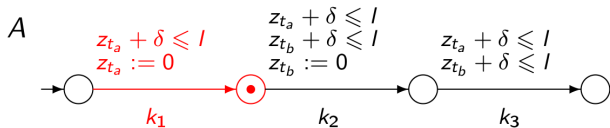
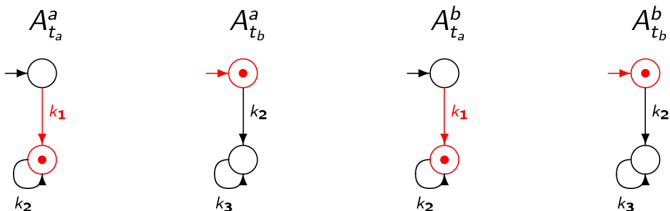
# Automata Model

Network of synchronized timed automata for NSPK protocol



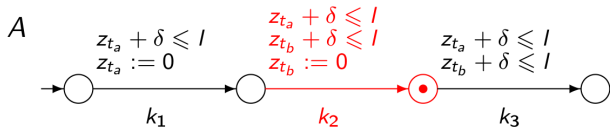
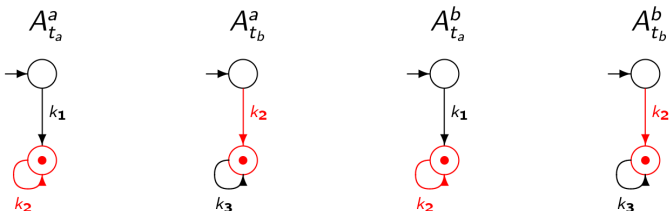
# Automata Model

Network of synchronized timed automata for NSPK protocol



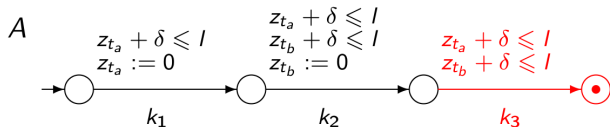
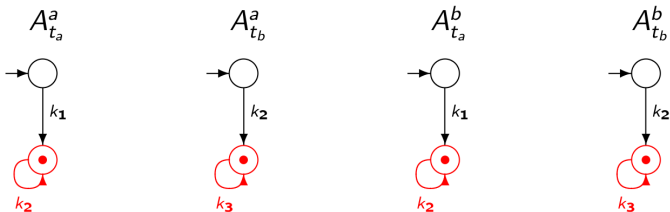
# Automata Model

## Network of synchronized timed automata for NSPK protocol



# Automata Model

## Network of synchronized timed automata for NSPK protocol



# Time dependencies

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## Message composing time

$$T_c = T_{sz} + T_g \quad (2)$$

## Duration of k-step

$$T_k = T_{sz} + T_g + D + T_{dsz} \quad (3)$$

$$T_k^{min} = T_{sz} + T_g + D_{min} + T_{dsz} \quad (4)$$

$$T_k^{max} = T_{sz} + T_g + D_{max} + T_{dsz} \quad (5)$$

# Time dependencies

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## Session time

$$T_{ses} = \sum_{k=1}^n T_k \quad (6)$$

$$T_{ses}^{min} = \sum_{k=1}^n T_k^{min} \quad (7)$$

$$T_{ses}^{max} = \sum_{k=1}^n T_k^{max} \quad (8)$$

# Time dependencies

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## Lifetime in single step

$$T_k^{out} = \sum_{i=k}^n T_i^{max} \quad (9)$$

where:

- $k$  – step number,
- $n$  – number of all steps in protocol,
- $T_i^{max}$  – maximum time of step execution.

# Steps of Procedure

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- Writing protocol in ProToc language.
- Generating a set of protocol executions.
- Generating a network of timed automata.
- Generating a formulas for SAT-Solver.
- SAT-Solver testing.
- Saving results to a file.



# Selected SAT-solvers

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Protocols	Mini-SAT	Lingeling	Clasp	Glucose	Treengeling
<b>Memory [MB]</b>					
NSPK	2,95	2,99	3,01	3,2	4
NSPK <sub>Low</sub>	3,95	4,4	4,02	4,23	5
<b>Time [ms]</b>					
NSPK	48	56	60	200	260
NSPK <sub>Low</sub>	360	400	370	560	430

# Experimental results for NSPK protocol

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Length of path	Variables	Clauses	Memory [MB]	Time [ms]	Result
2	6087	15781	2,19	8	UNSAT
3	7322	18829	2,19	8	UNSAT
4	13873	34298	2,48	20	UNSAT
5	15051	37188	2,55	28	UNSAT
6	21849	53246	2,95	48	SAT

Time assumptions:

- Delays in  $D = 0, 15[tu]$ ,
- lifetime  $Lf = 2[tu]$

# Susceptibility to Attack

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Protocol NSPK will be vulnerable to attack if

- in the first step:  $D \leq Lf$
- in the second step:  $D \leq Lf/4$
- in the last step:  $D \leq Lf/5$

# Experimental results for NSPK protocol, with time restrictions

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Length of path	Variables	Clauses	Memory [MB]	Time [ms]	Result
2	6216	16118	2,19	8	UNSAT
3	7533	19366	2,18	12	UNSAT
4	14177	35054	2,48	32	UNSAT
5	15431	38130	2,55	28	UNSAT
6	22334	54435	2,95	80	UNSAT

Time assumptions:

- Delay in the following steps  
 $D_1 = 10, 1[tu]$ ,  $D_2 = 2, 6[tu]$ ,  $D_3 = 2, 1[tu]$ ,
- lifetime  $Lf = 10[tu]$

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- Presented method can be used for fast and simple protocol verification.
- With the implemented tool, we can not only find the attack on the protocol, but also check if the protocol makes sense.
- Shown time constraints, enable us to determine the protocol time frame in which it is vulnerable to attack.
- This is one of the steps to accurately show the strengths and weaknesses of security protocols.

# References I

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